

- Compiler - source → target language	→ 86.64	- C/C++ compiler
- Interpreter - source → value	- res = return value	- raw numbers
- → CPU = 86.64	- rsp = stack pointer	- (add x), (sub x)
- Assembly → machine code	- rdi = our heap pointer	- true, false
- runtime - code base, stuff for compiler	→ mostly C	- (not e), (num? e), (seq? e)
- S-expression - lisp	- Hw1	- (if else <then cells>)
- () → ¹⁰⁰ of now	- desugaring	- (t + t), -, =, < for
- symbol, number, list	- Hw2	eval eval eval eval
- Unary expansion	- permitted characters	→ type checking, variable
- !ers	- !@#\$%^&*	→ local scope
- patternmatch	- newline	- (let ((var = el)) e2)
- Open Box character art	- space	- make variable only!
- testing	- character 0b 00001111	- (pair el e2), (left e), (right e)
- how worth it?	- Hw3	- (read-var)
- program	- short circuit	- (print e), (newline), (do el ... en)
- LISP compiler	- car	- (define (elmore const ...) e)
- formal verification	- cdr	- Hw4 - none
↳ language with functional interpreter	- tail-rec	- Hw5
- Boolean	- left-right	- (char? e), (char->num e), (num->char e)
- add type	- left-left-right	- Hw3
- switch symbols	- right	→ environment - and, or, multiple let, case
- "dynamic typing"	- pairchain	- *, /
- tag not tail 0x00	- C for RDI	- Hw4
- tag tail not 0xffffffff	- reaching 0b101	- list?, vector?, vector-length
- Conditionals	- Hw5	vector, vector-get vector-set
- false → only false value	- list	- Hw5
- x vs conditional jump, jz, cmp	- () entry	string > el111, stdin/stdout open, open-within-in, close-out
- else, continue labels	- pairchain	inputs, outputs
- cast from label ref to	- C for RDI	
- Binary operators	- align h phewels	
- need to "fix" one of them	- invariant, unchanged	
- streak	0b01111111 0b11111111	
- rsp = return address in stack	- C var	
- move down (rsp - 8)	var - RDI:	
- el at stack index, b to 8	2nd - RBP:	
- List	3rd - RDX:	
- symbols map var name to value	- Hw1	
- immutable, make new ones	- align r - RDI	
- when on 0x0000 to check	- Hw2	
- put on stack, symbol tracks it	- align r - RBP	
- Pairs	- Hw3	
- don't form stack on -> heap!	- Hw4	
- keep in RDI:	- Hw5	
- 64L addresses, multiply & + 0x0000	- mvw r, [rsp - 8]	
- no gc yet	- Drills	
- Errors	- mvw r, [rsp - 8]	
- call C function error		
- just jump to it, don't match tag		
- interpreter doesn't check "well-formed"		
- be decrement everything		
- User Input		
- OEM input fail msg exits code		
- call C function		
- Res rdi		
- move esp down (custom esp - ?)		
- returnable at esp, with call to main		
- align esp to 16		
- User Output		
- print return		
- C args		
- printing in RDI		
- Functions		
- sy add		
- don't use vars when calling		
- tail recursive		
- leaves esp - 8 to be returnable		
- 2nd esp - 16 → esp		
... last in stack?		
- Tail Call		
- reuse stack frame (don't & esp)		
- ↪ just jumps, no call		
- Parsing		
- parsing		
- from-U		

- First-class functions
 - function basically a value
 - anonymous function (closure)
 - closure = ref actual variable
 - interpreter
 - add taking holds for arguments
 - example:
 - same, closure address in use
- Lambda
 - anonymous function
 - doesn't have to be local and immutability means
- Closure
 - same as the environment
 - interpreter - just goes to env
 - compiler
 - makes you care "free variables"
 - Env ref?
 - add to stack, like closures
 - keep
 - pointer to begin of args
- Optimization
 - copy propagation analysis
 - global
 - Peephole
 - common subexpression elimination
 - dead code elimination
 - loop invariant code motion
 - better than (in. code)
 - what?
 - ASTs, variables
 - propagate (current)
 - order? chain pattern, phase ordering
 - visibility just one environment
 - Control Flow
 - recursive
 - check for field inheritance
 - Peephole optimization
 - look at small windows
 - else
 - algebraic simplification
 - else reduction and cleanup
 - else unreachable code
 - else loop invariant code
 - flow control optimization
 - J21 J22
J21 J22
J22 J21
- Garbage Collection
 - main idea
 - stop memory leak
 - but don't do it (but wouldn't know what)
 - In GC
 - mark and sweep
 - reference counting
 - mark and sweep
 - can't tell all environments
 - mark and sweep
 - for pointers and no copies
 - remove everything from heap and deallocate
 - AST for pointers
 - one slot for each cell to handle w/ free or
- cheap add & free
- Intermediate Representation (IR)
 - and ANF for between them also
- Register Allocation
 -
 - Method level
 - static, graph based search all reachable
 - propagates, adds variable to frontier
 - fragmentation
 - Block level (CPA)
 - 2 stages (need 2x memory)
 - copy unreachable, either 0 or 1 or both
 - Component CG
 - more efficient than graph level CG (no memoization)
 - Def Count CG
 - expensive
 - Register level allocation
 - 3. in registers